## Algebraic multigrid based on De Rham complexes on graphs with applications to Maxwell's equations

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17th International Conference on Domain Decomposition Methods St. Wolfgang/Strobl - July 3, 2006


## Formulation

Our aim is the resolution of linear systems coming from the edge finite element discretisation of:

$$
\left\{\begin{array}{l}
\text { To find } \mathbf{E} \in V \text { such that: } a\left(\mathbf{E}, \mathbf{E}^{\prime}\right)=F\left(\mathbf{E}^{\prime}\right), \forall \mathbf{E}^{\prime} \in V_{0}, \\
\text { with } a\left(\mathbf{E}, \mathbf{E}^{\prime}\right)=\int_{\Omega} \nu \operatorname{curl} \mathbf{E} \cdot \operatorname{curl} \mathbf{E}^{\prime}+\int_{\Omega} \gamma \mathbf{E} \cdot \mathbf{E}^{\prime}
\end{array}\right.
$$

with $V \subset \mathbf{H}($ curl,$\Omega)$. We consider $\nu$ and $\gamma>0$.

- Straightforward applications: transient vector wave equation or eddy current formulation.
- Possible extensions: magnetostatics, time-harmonic problems.


## Finite element spaces

Two finite element spaces considered :

- Edge element space $\mathcal{W}^{1}(\mathcal{T})$ included in $\mathbf{H}($ curl,$\Omega)$ : the lowest order Nedelec elements. Degrees of freedom: path integral along the mesh edges. Basis $\left(\lambda_{i}^{h}\right)_{i \in\left\{1, \ldots, E^{n}\right\}}$.
- Nodal element space $\mathcal{W}^{0}(\mathcal{T})$ included in $\mathrm{H}^{1}(\Omega)$ : $P_{1}$-Lagrange finite elements. Degrees of freedom: values at the mesh vertices. Basis $\left(\phi_{p}^{h}\right)_{p \in\left\{1, \ldots, N^{h}\right\}}$.

Coming from continuous space properties, both discrete spaces belong to a De Rham complex with complete sequence property. In particular:

$$
\begin{aligned}
& \qquad \operatorname{grad} \phi_{p}^{h}=\sum_{i=1}^{E^{h}} G_{i p}^{h} \lambda_{i}^{h}, \forall p \in\left\{1, \ldots, N^{h}\right\} . \\
& \text { and } \operatorname{range}(\mathrm{grad})=\operatorname{ker}(\text { curl }) .
\end{aligned}
$$

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$G^{h}=\left(\begin{array}{cccc}-1 & 1 & 0 & 0 \\ 0 & -1 & 0 & 1 \\ 0 & 0 & -1 & 1 \\ -1 & 0 & 1 & 0\end{array}\right)$

$$
\operatorname{grad} \phi_{p}^{h}=\sum_{i=1}^{E^{h}} G_{i p}^{h} \lambda_{i}^{h}, \forall p \in\left\{1, \ldots, N^{h}\right\}
$$

and range $($ grad $)=\operatorname{ker}($ curl $)$.

## Smoother choice

- Laplacian case: a few iterations of damped Jacobi or Gauss-Seidel provide an efficient smoothing.
- curl $\nu$ curl case: classical algorithms do not deal with the component in the kernel of curl, even if the dimension of this subspace is proportional to the number of mesh nodes.
- In order to supress this difficulty, a hybrid Gauss-Seidel algorithm has been proposed by Hiptmair ${ }^{1}$ :
- Classical smoothing for the global system.
- Specific smoothing on the kernel of curl.

[^0]
## Algebraic construction of prolongation operators

- The coarse nodal and edge spaces are included in the fine spaces:

$$
\begin{aligned}
\phi_{n}^{H} & =\sum_{p=1}^{N^{h}} P_{p n}^{\mathrm{nod}} \phi_{p}^{h}, \forall n \in\left\{1, \ldots, N^{H}\right\}, \\
\lambda_{e}^{H} & =\sum_{i=1}^{E^{h}} P_{i e}^{\mathrm{edg}} \lambda_{i}^{h}, \forall e \in\left\{1, \ldots, E^{H}\right\} .
\end{aligned}
$$

- In order to use the same smoother as in the geometric case, the following relation will be kept at the coarse level

$$
\begin{equation*}
\operatorname{grad} \phi_{n}^{H}=\sum_{e=1}^{E^{H}} G_{e n}^{H} \lambda_{e}^{H}, \forall n \in\left\{1, \ldots, N^{H}\right\} . \tag{1}
\end{equation*}
$$

$\left(\phi_{n}^{H}\right)_{n \in\left\{1, \ldots, N^{H}\right\}}$ is the coarse nodal basis and $\left(\lambda_{e}^{H}\right)_{e \in\left\{1, \ldots, E^{H}\right\}}$ is the coarse edge basis.

## Algebraic construction of prolongation operators

Relation (1) is implicitly satisfied in geometric multigrid.
This relation is wished to be kept for algebraic methods, in particular to define a coarse graph whose incident matrix $G^{H}$ will be used in this relation.
By gathering the previous relations, the commutativity relation is obtained:

$$
P^{\mathrm{edg}} G^{H}=G^{h} P^{\mathrm{nod}} .
$$

$G^{h}$ describes the fine graph, $P^{\text {nod }}$ is the nodal prolongator computed by existing methods.

To define or to compute:
$G^{H}$ and the edge prolongator $P^{\text {edg }}$.

## Remark on the computation of $P^{\text {nod }}$

Many algebraic multilevel methods for grad div operator can provide $P^{\text {nod }}$ operators satisfying the conditions required by our construction.
In particular, they build a decomposition of the domain $\Omega$ in overlapping subdomains $\left(\Omega_{n}^{H}\right)_{n=1, \ldots, N^{H}}$ such that:

$$
\operatorname{supp} \phi_{n}^{H} \subset \overline{\Omega_{n}^{H}}
$$

This is equivalent to enforce zero coefficients in the $n^{\text {th }}$ column of $P^{\text {nod }}$. It is also required that $P^{\text {nod }}$ satisfies a partition of unity property:

$$
\forall x \in \bar{\Omega}, \sum_{n=1}^{N^{H}} \phi_{n}^{H}(x)=1 \Leftrightarrow \sum_{n=1}^{N^{H}} P_{p n}^{\mathrm{nod}}=1, \forall p \in\left\{1, \ldots, N^{h}\right\} .
$$

## Coarse graph and support of the functions

After the computation of the nodal prolongator $P^{\text {nod: }}$

- Definition of a coarse graph or equivalently a coarse edge-node incidence matrix $G^{H}$.
For a coase edge of extremities coarse nodes $n$ and $m$, the condition $\Omega_{n}^{H} \cap \Omega_{m}^{H} \neq \emptyset$ has to be ensured.
- A coarse edge function $\lambda_{e}^{H}$ is associated with each coarse edge $e$ and $\operatorname{supp} \lambda_{e}^{H} \subset \overline{\Omega_{n}^{H} \cap \Omega_{m}^{H}}$.
As for $P^{\text {nod }}$, it is equivalent to enforce zero coefficients in the $e^{\text {th }}$ column of $P^{\text {edg }}$.
$I_{e}$ : Set of the non-zero coefficients in the $e^{\text {th }}$ column of $P^{\text {edg }}$.


## Illustration



## Illustration



## Final formulation

- Reitzinger and Schöberl ${ }^{2}$ proposed a strategy to build nodal and edge prolongators satisfying the relation $P^{\text {edg }} G^{H}=G^{h} P^{\text {nod }}$. Their approach is attached to a particular choice for the nodal prolongator $P^{\text {nod }}$ which leads to a unique matrix $P^{\text {edg }}$ satisfying the commutation property.
- We wished to extend this approach to other choices of the nodal prolongator $P^{\text {nod }}$.
Several matrices $P^{\text {edg }}$ can then satisfy the commutativity relation and supplementary rules should be defined to make a judicious choice.

[^1]
## Final formulation

Ideas coming from nodal finite element strategy: constraints on the supports (to keep the sparsity or the matrices) + minimisation of an energy functional + approximation constraints.

## Optimisation problem

$$
\left\{\begin{array}{l}
\text { Find } P^{\text {edg }} \text { minimising } \sum_{e=1}^{E^{H}} \beta_{e}^{t} K_{e} \beta_{e} \\
\text { under the constraint: } P^{\mathrm{edg}} G^{H}=G^{h} P^{\mathrm{nod}}
\end{array}\right.
$$

- $\beta_{e}$ : contains the non-zero coefficients of the $e^{\text {th }}$ column of $P^{\text {edg }}$, i.e. these indexed by $I_{e}$.
- $K_{e}$ : Local symmetric positive definite (SPD) matrix. The whole $K_{e}$ matrices define the energy functional for the minimisation.


## Subgraphs $\mathcal{S}^{H, i}$

$\mathcal{S}^{H, i}$ : graph defined for each fine edge $i$ such that one index $e$ satisfies $i \in I_{e}$.
Subgraph of the coarse graph obtained by keeping only the coarse edges of index $e$ such that $i \in I_{e}$.


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## Flow problems

- We have shown that solving $P^{\text {edg }} G^{H}=G^{h} P^{\text {nod }}$ is equivalent to solve on each subgraph $\mathcal{S}^{H, i}$ a flow problem:

$$
\left(G^{H, i}\right)^{t} P_{i}^{\mathrm{edg}}=\left(G^{h} P^{\mathrm{nod}}\right)_{i .} .
$$

- To determine: $P_{i}^{\text {edg }}$, vector containing the components of the $i^{\text {th }}$ row of $P^{\text {edg }}$ indexed by the indices of the edges of $\mathcal{S}^{H, i}$.
- Known: $\left(G^{h} P^{\text {nod }}\right)_{i}$, vector containing the components of the $i^{\text {th }}$ row of $G^{h} P^{\text {nod }}$ indexed by the indices of the nodes of $\mathcal{S}^{H, i}$.


## Resolution of flow problems

Solution of the flow problem:

$$
P_{i}^{\mathrm{edg}}=\left(P_{i}^{\mathrm{edg}}\right)^{\prime}+\left(P_{i}^{\mathrm{edg}}\right)^{\prime \prime}
$$

where:

- $\left(P_{i}^{\text {edg }}\right)^{\prime}$ is a particular solution of the flow problem, determined thanks to a spanning tree of the subgraph,
- $\left(P_{i}^{\text {edg }}\right)^{\prime \prime}$ belongs to the kernel of $\left(G^{H, i}\right)^{t}$ which is generated by the independent cycles of the subgraph.



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## Linear system for the minimisation

After determining the particular solutions, the minimisation problem is equivalent to the resolution of a linear system of the form:

$$
B^{t} D B \Gamma=-B^{t} D \widetilde{P^{\text {edg }}},
$$

where:

- $\Gamma$ is the vector whose components give the coefficients of $\left(\left(P_{i}^{\text {edg }}\right)^{\prime \prime}\right)^{t}$ in the bases of the kernel of $\left(G^{H, i}\right)^{t}$.
- the matrix $B$ gathers the basis vectors of these different kernels. It is a sparse full-rank matrix which is assembled during the resolution of flow problems.
- the matrix $D$ is block-diagonal and its diagonal blocks are the matrices $K_{e}$ involved in the energy functional.
- the vector $\widetilde{P}^{\text {edg }}$ gathers the particular solutions $\left(P_{i}^{\text {edg }}\right)^{\prime}$ from all flow problems.


## Properties of the linear system

- The matrices $K_{e}$ being SPD, $B^{t} D B$ is SPD and the system can be solved by the conjugate gradient method.
- Depending on the choice for $K_{e}$, the conditioning of $D$ may be strongly dependent of the global dimension of the problem.
- However, if $D=I d$, the conditioning of $B^{t} B$ is low and independent of the global dimension of the problem.

Moreover, a high accuracy is not required for the resolution of this system; this is only a setup phase before solving the initial system. The solution kept, whatever the accuracy is, always satisfies the commutativity relation.

## Definition of $\Omega_{n}^{H}$ and of $G^{H}$

- Partition of the nodes of the graph (initialy of the mesh):

$$
\left\{1, \ldots, N^{h}\right\}=\bigcup_{n=1}^{N^{H}} H_{n}
$$

Definition of subdomains $\left(\widetilde{\Omega}_{n}^{H}\right)_{n=1, \ldots, N^{H}}$ :

$$
\widetilde{\Omega}_{n}^{H}=\bigcup_{p \in H_{n}} \operatorname{supp}\left(\phi_{p}^{h}\right)
$$

- Creation of a coarse edge e of extremities $n, m$ iff $\widetilde{\Omega}{ }_{n}^{H} \cap \widetilde{\Omega}{ }_{m}^{H} \neq \emptyset$.
- Subdomain $\Omega_{n}^{H}$ defined by extending $\widetilde{\Omega}_{n}^{H}$ to the nearest neighbours. Without this extension, no degree of freedom for the minimisation: this is the Reitzinger and Schöberl method (RS method).


## Definition of $\Omega_{n}^{H}$ and of $G^{H}$



Initial graph.


Coarse graph representation.

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Initial graph.


Coarse graph representation.

## Choice for the matrices $K_{e}$

Different choices used for numerical experiments:

- Matrices extracted from the global matrix $K$ of the problem; it is denoted by K,
- Matrices extracted from the matrix defined from $\int_{\Omega} \nu$ curl $E \cdot$ curl $E^{\prime}$; it is denoted by $S_{\nu}$,
- Matrices are all equal to the identity; it is denoted by Id.


## 2D problem

- $\Omega=] 0 ; 1\left[^{2}\right.$.
- Tangential component of the field is enforced on the left side with $\mathbf{E}_{y}=\sin (2 \pi y)$. On the remaining of the boundary: $(\operatorname{curl} \mathbf{E}) \times \mathbf{n}=0$.
- Constant coefficient in the subdomains with the following equations:
- Basic test: curl curl $\mathbf{E}+\mathbf{E}=0$.
- Cavity in harmonic regime: curl curl $\mathbf{E}-\omega^{2} \mathbf{E}=0$ with $\omega=1.5 \pi$.

$\gamma=1$

$\gamma$ with $\omega=1.5 \pi$


## Mesh used



Initial mesh $\tau_{0}^{h}$


First refinement $\tau_{1}^{h}$

For the mesh $\tau_{i}^{h}, i$ indicates the number of refinement. NB: algebraic strategy works without knowledge of the grid structuration.

## Resolution of the optimisation linear system

|  | $\tau_{2}^{h}$ | $\tau_{3}^{h}$ | $\tau_{4}^{h}$ | $\tau_{5}^{h}$ | $\tau_{6}^{h}$ | $\tau_{7}^{h}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| optimisation system size | 62 | 253 | 1016 | 4081 | 16355 | 65479 |
| initial system size | 100 | 392 | 1552 | 6176 | 24640 | 98432 |

Table: Number of unknows for the optimisation linear system and for the initial linear system - 2D, structured meshes.

The number of unknowns for the minimisation is roughly $2 / 3$ of the number of unknowns for the initial system.

## Resolution of the optimisation linear system

Resolution by a conjugate gradient method without preconditioning.

|  | $\tau_{2}^{h}$ | $\tau_{3}^{h}$ | $\tau_{4}^{h}$ | $\tau_{5}^{h}$ | $\tau_{6}^{h}$ | $\tau_{7}^{h}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $K$ | 24 | 19 | 12 | 12 | 12 | 12 |
| $S_{\nu}$ | 14 | 12 | 12 | 12 | 12 | 12 |
| Id | 3 | 1 | 1 | 1 | 1 | 1 |

Table: Number of iterations for solving the optimisation linear system (division of the residual by $10^{3}$ ) -2 D , structured meshes.

- Results given only for the finest level.
- Number of iterations independent of the dimension of the global system.


## Resolution of the initial linear system - Case 1

- Conjugate gradient preconditioned by a multilevel method using Hiptmair's smoother.
- Stopping criterion: division by $10^{10}$ of the residual norm.
geom: geometric multigrid,
1 level: use only one level, smin: use only the particular solution of flow problems.



## Resolution of the initial linear system - Case 2

- General deterioration of the results for $\gamma=-(1.5 \pi)^{2}$.
- The hierarchy between the methods remains the same.



## Final results and perspectives

Final results

- Description of the edge prolongators satisfying the commutativity relation. Generic algorithms to compute these prolongators.
- Addition of an optimisation process in order to optimise the choice of the edge prolongator.
Perspectives
- Introduction of other methods for the construction of the coarse graph.
- Optimisation of basis satisfying the commutativity relation but initially computed by another strategy.


[^0]:    ${ }^{1}$ Multigrid method for Maxwell's equations. SIAM J. Numer. Anal., 36(1), 1999.

[^1]:    ${ }^{2}$ S. Reitzinger and J. Schöberl. An algebraic multigrid method for finite element discretizations with edge elements. Numer. Linear Algebra App., 9(3), 2002.

